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Algebraic Structure of Symbolic Expressions

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ABSTRACT

The ring S_∞ of formal power series in the noncommuting variables $a_1, \ldots, a_n$ with the coefficient field GF(2) is introduced and studied. The term *symbolic expressions* is used instead of formal power series, since it generalizes the concept of a symbolic expression introduced in [2] and [3]. S_∞ is characterized as the terminal object of the category **Aut** of automata. A category theoretic characterization of the subring S^{rat} of S_∞ consisting of rational sexps is also given.

Introduction

Theory of formal power series in noncommuting variables provides a useful algebraic tool for the study of formal languages. In the most general setting of the theory, the coefficients of a formal power series are taken from an arbitrary *semiring*; and it is possible to prove useful theorems in this general setting. (See e.g. Salomaa and Soittola[1].) In this paper, however, we will take the two elemented Galois field GF(2) as the coefficient semiring. This choice of the semiring will turn out to be very convenient. We will use the term *symbolic expression* (or *sexp* for short) instead of formal power series, since it generalizes the concept of a symbolic expression introduced in Sato[2] and Sato and Hagiya[3]. In [2] and [3], it is shown that symbolic expressions constitute a flexible data structure, which is therefore used as a data domain of a programming language called Hyperlisp. Here we study symbolic expressions from an algebraic point of view.

In 1, we study the ring S_∞ consisting of all the sexps. We show that S_∞ satisfies a certain domain equation for an abstract data structure.

In 2, we characterize S_∞ as the terminal object of the category **Aut** of automata. We then study the subring S^{rat} of S_∞ consisting of *rational* sexps. The well-known characterization of regular languages (i.e., rational sexps) in terms of finite automata is established in our formalism. A category theoretic characterization of S^{rat} is also given.

In 3, we introduce the subring S of S_∞ consisting of *finite* sexps. The relationship with the concept of a symbolic expression introduced in [2], [3] is also discussed.

1. S_∞

Let $\Sigma=\{a_1, \ldots, a_n\}$ $(n\geqq 1)$ be an alphabet consisting of n distinct symbols, which we will fix for the rest of this paper. Let $W=\Sigma^*$ be the free monoid over the alphabet Σ, and let $2=\{0,1\}$ be the Galois field GF(2). We put

$$S_\infty=\{r \mid r: W\to 2\}.$$

We will use r, s, t etc. to denote elements of S_∞ and u, v, w etc. to denote elements of W. Elements of S_∞ are called *symbolic expressions* or *sexps* for short. Elements of W are called *words*. For a word w, $|w|$ denotes its length. We write (r, w) for $r(w)$. We remark that S_∞ may be identified with 2^W (the power set of W) by the correspondence:

$$r \leftrightarrow \mathrm{supp}(r)=\{w\in W \mid (r, w)=1\}.$$

Any sexp r, then, naturally becomes a language over W.

We now define addition and multiplication on S_∞ as follows.

$$(r+s, w)=(r, w)+(s, w),$$
$$(rs, w)=\sum_{w=uv}(r, u)(s, v).$$

S_∞ then becomes a noncommutative ring with the 0 and 1 defined by

$$(0, w)=0,$$
$$(1, w)=\begin{cases} 1 & \text{if } w=1 \text{ (the unit of } W), \\ 0 & \text{otherwise.}\end{cases}$$

By identifying 0, $1\in S_\infty$ with those in 2, we assume that $2\subseteq S_\infty$. S_∞ then becomes a vector space over 2. Next, for any $w\in W$ we define $\bar{w}\in S_\infty$ by:

$$(\bar{w}, u)=\begin{cases} 1 & \text{if } w=u, \\ 0 & \text{otherwise.}\end{cases}$$

Since the map: $w\to\bar{w}$ is one-to-one and preserves multiplication on W, we will identify $\bar{w}$ with w and assume that $W\subseteq S_\infty$.

Consider the map $\pi: S_\infty\to S_\infty$ defined by

$$\pi(r)=(r, 1).$$

It is a ring homomorphism and satisfies $\pi^2=\pi$. If we regard S_∞ as a vector space, π becomes a projection and we have the direct sum decomposition:

$$S_\infty=\mathrm{Im}\,\pi\oplus\mathrm{Ker}\,\pi.$$

Since $\mathrm{Im}\,\pi = 2$, if we put

$$\mathbf{M}_\infty = \mathrm{Ker}\,\pi = \{r \mid (r, 1) = 0\}$$

we have

$$\mathrm{S}_\infty = 2 \oplus \mathbf{M}_\infty \tag{1.1}$$

We put $\mathbf{A}_\infty = \mathrm{S}_\infty - \mathbf{M}_\infty$. Elements of $\mathbf{M}_\infty$ are called *molecules*, and elements of $\mathbf{A}_\infty$ are called *atoms*.

Next, we define a map $\delta : \mathrm{S}_\infty \times W \to \mathrm{S}_\infty$ by:

$$(\delta(r, u), v) = (r, uv).$$

It is an action of the monoid W on S_∞, since we have

$$\delta(r, 1) = r,$$
$$\delta(\delta(r, u), v) = \delta(r, uv).$$

For a fixed w,

$$\delta(-, w) : \mathrm{S}_\infty \to \mathrm{S}_\infty$$

is a linear transformation. In particular, for each i $(1 \leqq i \leqq n)$, we define $\sigma_i : \mathrm{S}_\infty \to \mathrm{S}_\infty$ by

$$\sigma_i(r) = \delta(r, a_i).$$

We have the following

Proposition 1.1. $\sigma_i(st) = \pi(s)\sigma_i(t) + \sigma_i(s)t \quad (1 \leqq i \leqq n)$.

Proof.

$$\begin{aligned}
(\sigma_i(st), w) &= (\delta(st, a_i), w) \\
&= (st, a_i w) \\
&= \sum_{uv = a_i w} (s, u)(t, v) \\
&= (s, 1)(t, a_i w) + \sum_{a_i uv = a_i w} (s, a_i u)(t, v) \\
&= \pi(s)(\sigma_i(t), w) + \sum_{uv = w} (\sigma_i(s), u)(t, v) \\
&= (\pi(s)\sigma_i(t) + \sigma_i(s)t, w).
\end{aligned}$$

Note that, by a simple computation, we have $\sigma_i(a_j) = \delta_{ij}$, where δ_{ij} is Kronecker's delta.

We now regard S_∞ as a right S_∞-module. $\mathbf{M}_\infty$ then becomes its submodule (or, a right ideal of the ring S_∞).

Theorem 1.2. $< a_1, \ldots, a_n >$ *forms a basis of* $\mathbf{M}_\infty$.

Proof. It suffices to prove the following (a) and (b).

(a) If $r\in \mathbf{M}_\infty$ then $r=\sum_i a_i\sigma_i(r)$: Since $r\in\mathbf{M}_\infty$, we have $(r,1)=0$. On the other hand, since $a_i\in\mathbf{M}_\infty$,

$$(\sum_i a_i\sigma_i(r),1)=\sum_i (a_i\sigma_i(r),1)=\sum_i (a_i,1)(\sigma_i(r),1)=0.$$

Next, for any $w\in W$ we have

$$\begin{aligned}(\sum_i a_i\sigma_i(r),a_jw)&=(\sigma_j(\sum_i a_i\sigma_i(r)),w)\\&=(\sum_i \sigma_j(a_i\sigma_i(r)),w)\\&=\sum_i(\sigma_j(a_i\sigma_i(r)),w)\\&=\sum_i[(\sigma_j(a_i)\sigma_i(r),w)+(\pi(a_i)\sigma_j(\sigma_i(r)),w)]\\&=\sum_i(\delta_{ji}\sigma_i(r),w)\\&=(\sigma_j(r),w)\\&=(r,a_jw).\end{aligned}$$

Since any $u\in W$ is either 1 or of the form a_jw we have $r=\sum_i a_i\sigma_i(r)$.

(b) $\sum_i a_it_i=0 \implies t_i=0\ (1\leqq i\leqq n)$:

$$0=\sigma_j(\sum_i a_it_i)=\sum_i\sigma_j(a_it_1)=\sum_i\sigma_j(a_i)t_i+\pi(a_i)\sigma_j(t_i)=\sum_i\delta_{ji}t_i=t_i.$$

By this theorem we have the right $\mathbf{S}_\infty$-module isomorphism:

$$\sigma:\mathbf{S}_\infty\oplus\cdots\oplus\mathbf{S}_\infty\to\mathbf{M}_\infty \tag{1.2}$$

such that $\sigma(t_1,\ldots,t_n)=a_1t_1+\cdots+a_nt_n$. We have

$$\sigma^{-1}(r)=<\sigma_1(r),\ldots,\sigma_n(r)>.$$

In view of (1.1), the map

$$\tau:\mathbf{S}_\infty\times\cdots\times\mathbf{S}_\infty\to\mathbf{A}_\infty$$

defined by

$$\tau(t_1,\ldots,t_n)=\sigma(t_1,\ldots,t_n)+1$$

is a bijection. Combining (1.1) and (1.2), we have the following set theoretic isomorphism:

$$\mathbf{S}_\infty\simeq 2\times\mathbf{S}_\infty\times\cdots\times\mathbf{S}_\infty \tag{1.3}$$

where $r\leftrightarrow<\pi(r),\sigma_1(r),\ldots,\sigma_n(r)>$. We have the following proposition by

(1.3).

Proposition 1.3. $s=t \iff \pi(s)=\pi(t), \sigma_i(s)=\sigma_i(t) \quad (1 \leqq i \leqq n)$.

(1.3) may be rewritten as:

$$\mathbf{S}_\infty \simeq \mathbf{S}_\infty^n + \mathbf{S}_\infty^n \tag{1.4}$$

where + denotes the (direct) sum of two sets. This isomorphism tells us the basic properties of the data structure $\mathbf{S}_\infty$. Namely, any sexp is an infinite n-ary tree which carries one bit information at each node. The *recognizer* π distinguishes atoms from molecules. The *constructor* σ (τ) constructs from given n sexps t_i $(1 \leqq i \leqq n)$ a molecule (atom, resp.) s whose i-th subtree t_i is recovered by the *selector* σ_i.

A sexp s is *invertible* if there exists a t such that $st=ts=1$. Since the t above is unique for an invertible s, it is called the *inverse* of s and is denoted by s^{-1}. We wish to characterize invertible sexps. We need the following lemma.

Lemma 1.4. *If* $r \in \mathbf{M}_\infty$ *then* $r^k \in M_k$ $(k \geqq 0)$ *where*

$$M_k = \{ r \in \mathbf{S}_\infty \mid |w| < k \Longrightarrow (r, w) = 0 \}.$$

Proof. If $k=0$ then $r^0=1 \in M_0$. Assume $r^k \in M_k$. Then for any w such that $|w| < k+1$,

$$\begin{aligned}(r^{k+1}, w) &= \sum_{uv=w} (r^k, u)(r, v) \\ &= (r^k, w)(r, 1) + \sum_{\substack{uv=w \\ u \neq w}} (r^k, u)(r, v).\end{aligned}$$

Since $r \in \mathbf{M}_\infty$, we have $(r, 1)=0$; and if $u \neq w$ we have $(r^k, u)=0$ by the assumption. Hence $(r^{k+1}, w)=0$.

Theorem 1.5. *A sexp is invertible iff it is an atom.*

Proof. $(\Longrightarrow)$ If s is invertible, then $ss^{-1}=1$. Hence, $1= \pi(1)= \pi(ss^{-1})=\pi(s)\pi(s^{-1})$. Then we have $\pi(s)=1$, so s is an atom.

$(\Longleftarrow)$ Let s be an atom. We define a molecule r by putting $r=1+s$. Then we define a sexp t by:

$$(t, w) = (1 + r + \cdots + r^{|w|}, w).$$

By Lemma 1.4, for any $k \geqq 0$, we have

$$(1 + r + \cdots + r^{|w|+k}, w) = (1 + r + \cdots + r^{|w|}).$$

We have $st=1$ because:

$$\begin{aligned}(st, w) &= \sum_{uv=w} (s, u)(t, v) \\ &= \sum_{uv=w} (1+r, u)(1 + r + \cdots + r^{|v|}, v) \\ &= \sum_{uv=w} (1+r, u)(1 + r + \cdots + r^{|w|}, w)\end{aligned}$$

$$=(1+r^{|w|+1},w)$$

$$=(1,w)+(r^{|w|+1},w)$$

$$=(1,w).$$

That $ts=1$ holds can be proved similarly.

2. S^{rat}

We define S^{rat} as the least subset of S_∞ such that

(i) $2\cup\Sigma\subseteq\mathrm{S}^{rat}$,

(ii) $s,t\in\mathrm{S}^{rat}\Longrightarrow s+t\in\mathrm{S}^{rat}$,

(iii) $s,t\in\mathrm{S}^{rat}\Longrightarrow st\in\mathrm{S}^{rat}$,

(iv) $s\in\mathrm{S}^{rat}\cap\mathbf{M}_\infty\Longrightarrow(1+s)^{-1}\in\mathrm{S}^{rat}$.

S^{rat} is a subring of S_∞. In this section, we will study the relationship between S^{rat} and finite automata. Here we define an *automaton* (over Σ) as a triple

$$X=<X;\delta_X,\epsilon_X>$$

where

(1) X is a (possibly infinite) nonempty set of *states*,

(2) $\delta_X:X\times W\to X$ is an action of W on X,

(3) $\epsilon_X:X\to 2$.

Let X be an automaton. For each i ($1\leqq i\leqq n$), we define the map

$$\sigma_i^X:X\to X$$

by putting $\sigma_i^X(x)=\delta_X(x,a_i)$. This function determines the transition of states for the input symbol a_i. A state $x\in X$ is considered to be *accepted* if $\epsilon_X(x)=1$. We now define a function

$$L_X:X\to\mathrm{S}_\infty$$

by putting $(L_X(x),w)=\epsilon_X(\delta_X(x,w))$. $L_X(x)$ may be considered as the language which X, with the initial state x, accepts. Here we also note that $\mathrm{S}_\infty=<\mathrm{S}_\infty;\delta,\pi>$ becomes an automaton. Moreover, L_X becomes a morphism in the category **Aut** of automata which we now define.

The category **Aut**, by definition, has all automata as its objects. Its morphisms are defined by:

$h\in\mathrm{Hom}(X,Y)\Longleftrightarrow h$ is a map for which the diagram below commutes:

$$\begin{tikzcd}
X \times W \arrow[r, "h \times 1_W"] \arrow[d, "\delta_X"'] & Y \times W \arrow[d, "\delta_Y"] \\
X \arrow[r, "h"] \arrow[d, "\epsilon_X"'] & Y \arrow[d, "\epsilon_Y"] \\
2 \arrow[r, "1_2"'] & 2
\end{tikzcd}$$

Proposition 2.1. $L_X : X \to \mathrm{S}_\infty \in \mathrm{Hom}(X, \mathrm{S}_\infty)$.

Proof.

$$(\delta(L(x), w), u) = (L(x), wu) = \epsilon(\delta(x, wu)),$$

$$(L(\delta(x, w)), u) = \epsilon(\delta(\delta(x, w), u)) = \epsilon(\delta(x, wu)).$$

$$\pi(L(x)) = (L(x), 1) = \epsilon(\delta(x, 1)) = \epsilon(x).$$

Proposition 2.2. $L_{\mathrm{S}_\infty} : \mathrm{S}_\infty \to \mathrm{S}_\infty$ *is identity.*

Proof.

$$(L(r), w) = \pi(\delta(r, w)) = (\delta(r, w), 1) = (r, w1) = (r, w).$$

Proposition 2.3. $h \in \mathrm{Hom}(X, Y) \Longrightarrow L_Y \circ h = L_X$.

Proof.

$$(L_Y(h(x)), w) = \epsilon_Y(\delta_Y(h(x), w)) = \epsilon_Y(h(\delta_X(x, w))) = \epsilon_X(\delta_X(x, w)) = (L_X(x), w).$$

These propositions yield the following theorem.

Theorem 2.4. S^{rat} *is the terminal object of* **Aut**.

Proof. Let X be any automaton. We have $L_X \in \mathrm{Hom}(X, \mathrm{S}_\infty)$ by Proposition 2.1. Next, take any $h \in \mathrm{Hom}(X, \mathrm{S}_\infty)$. By Proposition 2.2 and Proposition 2.3 for $Y = \mathrm{S}_\infty$, we have $L_X = L_{\mathrm{S}_\infty} \circ h = 1 \circ h = h$. Thus we have proved that $\mathrm{Hom}(X, \mathrm{S}_\infty)$ is a singleton set for any X, i.e., S_∞ is terminal in **Aut**.

We now wish to characterize S^{rat} categorically. For a ring R, we let $M_k(R)$ denote the matrix ring consisting of $k \times k$ R-matrix. We define a ring homomorphism

$$\Pi_k : M_k(\mathrm{S}_\infty) \to M_k(2)$$

by putting $\Pi_k(S) = (\pi(s_{ij}))$ for $S = (s_{ij}) \in M_k(\mathrm{S}_\infty)$. The set

$$G_k = \Pi_k^{-1}(I),$$

where I is the $k \times k$ unit matrix, then becomes a monoid under matrix multiplication. Moreover, we have:

Theorem 2.5. *G_k forms a group under matrix multiplication.*

Proof. Let E_{ij} $(1 \leqq i, j \leqq k)$ be the $k \times k$ matrix such that its (i, j) element is 1 and all other elements are 0. For any molecule $r \in \mathbf{M}_\infty$, we put

$$Q_k(i, j; r) = I + rE_{ij}.$$

It is easy to see that $Q_k(i, j; r) \in G_k$ and

$$Q_k(i, j; r)^{-1} = \begin{cases} Q_k(i, j; 1 + (1+r)^{-1}) & \text{if } i = j, \\ Q_k(i, j; r) & \text{if } i \neq j. \end{cases}$$

We can then prove, using usual sweep out method, that the group generated by the set $\{Q_k(i, j; r)\}$ coincides with G_k.

Remark. The proof also shows that if $S \in G_k \cap M_k(\mathrm{S}^{rat})$, S^{-1} is also a member of $M_k(\mathrm{S}^{rat})$.

Let $X = \langle X; \delta_X, \epsilon_X \rangle$ be a finite automaton with k states so that $X = \{x_1, \ldots, x_k\}$. For each l $(1 \leqq l \leqq n)$ we define $\tilde{\sigma}_l : \{1, \ldots, k\} \to \{1, \ldots, k\}$ by the condition:

$$\tilde{\sigma}_l(i) = j \Longleftrightarrow \sigma_l^X(x_i) = x_j.$$

We then define a $k \times k$ $\mathbf{S}_\infty$-matrix $S = (s_{ij})$ by putting

$$s_{ij} = \delta_{ij} + \sum_l a_l \delta_{\tilde{\sigma}_l(i) j}.$$

We note that $S \in G_k \cap M_k(\mathrm{S}^{rat})$. Let X_i $(1 \leqq i \leqq k)$ be k distinct indeterminates and let $\mathbf{x} = {}^t(X_1 \cdots X_k)$. We also put $\mathbf{e} = {}^t(\epsilon(x_1) \cdots \epsilon(x_k))$. We call the equation:

$$S\mathbf{x} = \mathbf{e} \tag{2.1}$$

the *characteristic equation* of the finite automaton X. By Theorem 2.5 it has the unique solution $\mathbf{x} = S^{-1}\mathbf{e}$. Remark that (2.1) is equivalent to the following system of equations:

$$X_i = a_1 X_{\tilde{\sigma}_1(i)} + \cdots + a_n X_{\tilde{\sigma}_n(i)} + \epsilon(x_i) \quad (i = 1, \ldots, k).$$

Theorem 2.6.

$$L(x_i) = a_1 L(x_{\tilde{\sigma}_1(i)}) + \cdots + a_n L(x_{\tilde{\sigma}_n(i)}) + \epsilon(x_i) \quad (i = 1, \ldots, k).$$

Proof. Since $L \in \mathrm{Hom}(X, \mathrm{S}_\infty)$ we have

$$\sigma_l(L(x_i)) = L(\sigma_l^X(x_i)) = L(x_{\tilde{\sigma}(i)}),$$

$$\pi(L(x_i)) = \epsilon(x_i).$$

Therefore we have:

$$\pi(\mathrm{RHS}) = \pi(a_1 L(x_{\tilde{\sigma}_1(i)}) + \cdots + \pi(a_k L(x_{\tilde{\sigma}_k(i)}) + \pi(\epsilon(x_i))$$

$$
\begin{aligned}
&=\epsilon(x_i)\\
&=\pi(\mathrm{LHS}),\\
\sigma_l(\mathrm{RHS})&=\sigma_l(a_1 L(x_{\bar\sigma_1(i)}))+\cdots+\sigma_l(a_k L(x_{\bar\sigma_k(i)}))+\sigma_l(\epsilon(x_i))\\
&=L(x_{\bar\sigma_l(i)})\\
&=\sigma_l(L(x_i))\\
&=\sigma_l(\mathrm{LHS}).
\end{aligned}
$$

This proves LHS=RHS.

This theorem says that $L(x_i)$'s give the solution to the equation (2.1) and hence they are in $\mathbf{S}^{rat}$.

We next show that, conversely, any language in $\mathbf{S}^{rat}$ can be represented by a finite automaton. First we remark that, for a finite set X, 2^X becomes a vector space over 2 under the addition defined by:

$$U+V=(U-V)\cup(V-U).$$

If we identify any $x\in X$ with the singleton set $\{x\}$ then X becomes a basis of the vector space 2^X. Let $\mathbf{V}$ be any vector space over 2. Then any map $f: X\to\mathbf{V}$ can be uniquely extended to a linear map from 2^X to $\mathbf{V}$. We will denote this extended map also by f. We will write

$$X\ni x\models r$$

if x is a state of a finite automaton X and $r=L_X(x)$; and in this case we say that $x\in X$ *realizes* r. Such r's are called *realizable*.

Theorem 2.7. *A sexp r is realizable iff $r\in\mathbf{S}^{rat}$.*

Proof. Only if part follows from the remark after Theorem 2.4.

We prove *if part* by induction on the construction of r.

(i) Since the set $2\cup\Sigma\subseteq\mathbf{S}_\infty$ is closed under the functions σ_l $(1\leqq l\leqq n)$, it naturally becomes a finite automaton and each state realizes itself.

(ii) $r=s+t$: Assume that $X\ni x_0\models s$ and $Y\ni y_0\models t$. We define an automaton Z by putting:

$$Z=X\times Y=\{x\times y \mid x\in X, y\in Y\},$$

$$\delta_Z(x\times y, w)=\delta_X(x,w)\times\delta_Y(y,w),$$

$$\epsilon_Z(x\times y)=\epsilon_X(x)+\epsilon_Y(y).$$

Then by a simple computation we have $L_Z(x\times y)=L_X(x)+L_Y(y)$, so that $Z\ni x_0\times y_0\models s+t$.

(iii) $r=st$: Assume that $X\ni x_0\models s$ and $Y\ni y_0\models t$. We define an automaton Z by putting:

$$Z=2^Y\times X=\{\mathbf{y}\times x \mid \mathbf{y}\in 2^Y, x\in X\},$$

$$\sigma_l^Z(\mathbf{y}\times x)=(\sigma_l^Y(\mathbf{y})+\epsilon_X(x)\sigma_l^Y(y_0))\times\sigma_l^X(x)\quad(1\leqq l\leqq n),$$
$$\epsilon_Z(\mathbf{y}\times x)=\epsilon_Y(\mathbf{y})+\epsilon_X(x)\epsilon_Y(y_0).$$

We show that

$$\tilde{L}(\mathbf{y}\times x)=L_Y(\mathbf{y})+L_X(x)L_Y(y_0)\quad(\mathbf{y}\times x\in Z)$$

solves the characteristic equation of the automaton Z. I.e., we show that

$$\tilde{L}(z)=a_1\tilde{L}(\sigma_1^Z(z))+\cdots+a_n\tilde{L}(\sigma_n^Z(z))+\epsilon_Z(z)\quad(z\in Z).\tag{2.2}$$

Letting $z=\mathbf{y}\times x$, we compare the LHS and RHS of (2.2) as follows.

$$\begin{aligned}\pi(\mathrm{LHS})&=\pi(L_Y(\mathbf{y}))+\pi(L_X(x))\pi(L_Y(y_0))\\&=\epsilon_Y(\mathbf{y})+\epsilon_X(x)\epsilon_Y(y_0)\\&=\epsilon_Z(z)\\&=\pi(\mathrm{RHS}),\end{aligned}$$

$$\begin{aligned}\sigma_l(\mathrm{LHS})&=\sigma_l(L_Y(\mathbf{y}))+\sigma_l(L_X(x)L_Y(y_0))\\&=L(\sigma_l^Y(\mathbf{y}))+\pi(L_{X(x))}\sigma_l(L_X(y_0))+\sigma_l(L_X(x))L_Y(y_0)\\&=L(\sigma_l^Y(\mathbf{y}))+\epsilon_X(x)L_Y(\sigma_l^Y(y_0))+L_X(\sigma_l^X(x))L_Y(y_0)\\&=\tilde{L}(\sigma_l^Z(z))\\&=\sigma_l^Z(\mathrm{RHS})\quad(1\leqq l\leqq n).\end{aligned}$$

This proves (2.2), so that we have $L_Z(z)=\tilde{L}(z)$. Hence, we have $L_Z(\phi\times x_0)=L_X(x_0)L_Y(y_0)=st$. I.e., $Z\ni\phi\times x_0\models st$.

(iv) $r=(1+s)^{-1}$, $s\in\mathbf{S}^{rat}\cap\mathbf{M}_\infty$: Assume that $X\ni x_0\models s$. Since $1=r^{-1}r=(1+s)r=r+sr$, we have $r=1+sr$. So, $\sigma_l(r)=\pi(s)\sigma_l(r)+\sigma_l(s)r=\sigma_l(s)r$ $(1\leqq l\leqq n)$. We define an automaton Z by putting

$$Z=2^X=\{\mathbf{x}\mid\mathbf{x}\subseteq X\},$$
$$\sigma_l^Z(\mathbf{x})=\sigma_l^X(\mathbf{X})+\epsilon_X(\mathbf{x})\sigma_l^X(x_)\quad(1\leqq l\leqq n),$$
$$\epsilon_Z(\mathbf{x})=\epsilon_X(\mathbf{x}).$$

We show that

$$\tilde{L}(\mathbf{x})=L_X(\mathbf{x})r\quad(\mathbf{x}\in Z)$$

solves the characteristic equation of the automaton Z. I.e., we show the equation:

$$\tilde{L}(\mathbf{x})=a_1\tilde{L}(\sigma_1^Z(\mathbf{x}))+\cdots+a_n\tilde{L}(\sigma_n^Z(\mathbf{x}))+\epsilon_Z(\mathbf{x})\quad(\mathbf{x}\in Z)\tag{2.3}$$

We compare the LHS and RHS of (2.3) as follows.

$$\pi(\mathrm{LHS})=\pi(L_X(\mathbf{x}))\pi(r)=\epsilon_X(\mathbf{x})=\epsilon_Z(\mathbf{x})=\pi(\mathrm{RHS}),$$
$$\sigma_l(\mathrm{LHS})=\sigma_l(L_l^X(\mathbf{x})r)$$

$$
\begin{aligned}
&= \sigma_l(L_X(\mathbf{x}))r + \pi(L_X(\mathbf{x}))\sigma_l(r) \\
&= L_X(\sigma_l^X(\mathbf{x}))r + \epsilon_X(\mathbf{x})\sigma_l(s)r \\
&= (L_X(\sigma_l^X(\mathbf{x})) + \epsilon_X(\mathbf{x})L_X(\sigma_l^X(x_0)))r \\
&= L_X(\sigma_l^X(\mathbf{x}) + \epsilon_X(\mathbf{x})\sigma_l^X(x_0))r \\
&= \tilde{L}(\sigma_l^X(\mathbf{x}) + \epsilon_X(\mathbf{x})\sigma_l^X(x_0)) \\
&= \tilde{L}(\sigma_l^Z(\mathbf{x})) \\
&= \sigma_l(\mathrm{RHS}).
\end{aligned}
$$

This proves (2.3), so that we have $L_Z(\mathbf{x}) = \tilde{L}(\mathbf{x})$. Hence we have $Z \ni x_0 \models \tilde{L}(x_0) = L_X(x_0)r = sr = 1 + r$. By (ii) above, r is also realizable.

Corollary 2.8. $r \in \mathbf{S}^{rat} \Longrightarrow \sigma_l(r) \in \mathbf{S}^{rat}$ $(1 \leqq l \leqq n)$.

Proof. By Theorem 2.7, we can find X and x such that $X \ni x \models r$. Then we have $X \ni \sigma_l^X(x) \models \sigma_l(r)$. Hence $\sigma_l(r) \in \mathbf{S}^{rat}$.

Remark. It is possible to prove Corollary 2.8 directly by induction on the construction of r.

Let X be any automaton and let Y be a subset of X which is closed under σ_l^X for each l $(1 \leqq l \leqq n)$. Then we can naturally introduce into Y an automaton structure, by restricting that of X to Y, which makes Y a *subautomaton* of X. Since $\mathbf{S}^{rat}$ is closed under σ_l for each l $(1 \leqq l \leqq n)$, $\mathbf{S}^{rat}$ becomes a subautomanton of $\mathbf{S}_\infty$. Although $\mathbf{S}^{rat}$ is not a finite automaton, it is a *locally* finite automaton in the sense of the following definition.

Definition. An automaton $X = \langle X; \delta, \epsilon \rangle$ is *locally finite* iff the set $X|x = \{y \mid y = \delta(x, w)$ for some $w \in W\}$ is finite for all $x \in X$.

We will denote by $\mathbf{Aut}^{rat}$ the full subcategory of $\mathbf{Aut}$ consisting of all the locally finite automata. We have the following theorem.

Theorem 2.9. $\mathbf{S}^{rat}$ *is the terminal object of* $\mathbf{Aut}^{rat}$.

Proof. We first prove the claim that $\mathbf{S}^{rat}$ is a locally finite automaton. Suppose that $r \in \mathbf{S}^{rat}$. By Theorem 2.7, we can find X and x such that $X \ni x \models r$. Since $\mathrm{Im}\, L_X$ is finite and closed under σ_i $(1 \leqq i \leqq n)$, the set $X|x$ is also finite. This proves the claim.

Next, let X be an arbitrary locally finite automaton, and consider the map $L_X : X \to \mathbf{S}_\infty$. For any $x \in X$, $X|x$ becomes a finite subautomaton of X. Then we have $L_X(x) = L_{X|x}(x) \in \mathbf{S}^{rat}$, so that we may regard L_X as the map $L_X : X \to \mathbf{S}^{rat}$. Now the theorem can be proved similarly as Theorem 2.4.

3. S

We define $\mathbf{S}$ as the least subset of $\mathbf{S}_\infty$ such that

(i) $2 \cup \Sigma \subseteq S$,

(ii) $s, t \in S \Longrightarrow s+t \in S$,

(iii) $s, t \in S \Longrightarrow st \in S$.

According to this definition of S, S becomes a subring of S_∞. Elements of S are called *finite* sexps. We can establish the set theoretic isomorphism:

$$S \simeq S^n + S^n \tag{3.1}$$

similarly as (1.4). Just as (1.4) expressed some characteristics of S_∞, this equation says that S is a data structure equipped with the recognizer π, constructors σ, τ and selectors σ_i $(1 \leqq i \leqq n)$. Furthermore, it is easy to verity that S can be characterized as the least subset of S_∞ such that

(1) $0 \in S$,

(2) $t_1, \ldots, t_n \in S \Longrightarrow \sigma(t_1, \ldots, t_n) \in S$,

(3) $t_1, \ldots, t_n \in S \Longrightarrow \tau(t_1, \ldots, t_n) \in S$.

We remark that Scott[4] (p. 96) also discusses the domain equation of the form (3.1), and gives a solution for it as a *neighbourhood system*. In Scott[4], the interpretations of sums and products are slightly different, so that *total* elements in his solution corresponds to symbolic expressions in our sense. He also points out that *eventually periodic* total trees (which correspond to our rational sexp) represents an automaton which accepts itself.

Finally, we remark that in case $\Sigma = \{a_1, a_2\}$ finite sexps are precisely the symbolic expressions in the sense of Sato[2] and Sato and Hagiya[3]. In [2] and [3], the functions σ, τ, σ_1 and σ_2 are respectively called *cons*, *snoc*, *car* and *cdr* following the tradition of Lisp.

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